ITRI-WiMAXT: A WiMAX conformance testing tool

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Summary

The WiMAX Forum certifies and promotes the compatibility and interoperability of broadband wireless products where many testcases need to be developed for WiMAX procedures. In this paper, we develop a conformance test tool called ITRI-WiMAXT based on the TTCN-3 specifications. Then we show how WiMAX procedures are tested in ITRI-WiMAXT. Based on ITRI-WiMAXT, we have proposed several test procedures accepted by the WiMAX Forum. Copyright © 2009 John Wiley & Sons, Ltd.

KEY WORDS: conformance test; testing and test control notation version 3 (TTCN-3); worldwide interoperability for microwave access (WiMAX)

1. Introduction

Worldwide interoperability for microwave access (WiMAX) is a broadband wireless system based on the IEEE 802.16 standard, which supports high-speed wireless access to provide fixed and mobile broadband services over large coverage areas [1,2].

A simplified WiMAX network reference model is shown in Figure 1. The model is composed of three parts: mobile stations (MS; Figure 1 (1)), access service network (ASN; Figure 1 (2)) and the connectivity service network (CSN; Figure 1 (3)). The ASN comprises base stations (BSs; Figure 1 (4)) and ASN gateways (ASN-GW; Figure 1 (5)). A BS communicates with the MSs through radio links. Each BS is connected to an ASN-GW through the R6 interface. The BSs also communicate with each other for handover of data connections. The ASN-GW performs user data management for an MS through the BS. The CSN connects to the ASN to offer internet protocol (IP) connectivity. The entities residing in the CSN include the AAA server (Figure 1 (6)) that performs authentication, authorization and accounting, the home agent (HA; Figure 1 (7)) that provides mobile IP functionality [3] and the DHCP server (Figure 1 (8)) that allocates the IP addresses.

The WiMAX Forum certifies and promotes the compatibility and interoperability of broadband wireless products where many testcases need to be developed for WiMAX procedures. With the support of National Telecommunications Program (NTP) and Industrial Technology Research Institute (ITRI), we have proposed several test procedures [4,5] accepted by the WiMAX forum.

Based on the Testing and Test Control Notation version 3 (TTCN-3), this paper develops a conformance test tool called ITRI-WiMAXT. TTCN-3 is a high-level testing language used by European Telecommunication Standards Institute (ETSI) and WiMAX Forum.

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to create publicly available standardized test specifications. It precisely describes the test behaviour by defining the meaning of ‘pass’ and ‘fail’ for a testcase. This paper will show how WiMAX R6 message flows are tested in ITRI-WiMAXT. Note that based on the same tool, we have also developed conformance test tools for open mobile alliance (OMA) push-to-talk [6], download service [7], multimedia messaging service, instant message and presence service and browsing [8].

The paper is organized as follows. Section 2 describes the WiMAX message format and the initial network entry procedure. Section 3 introduces the TTCN-3 test system. Section 4 proposes the ITRI-WiMAXT test system. Section 5 provides our conclusions.

2. WiMAX Message Format and Initial Network Entry Procedure

The format of the WiMAX messages delivered in interface R6 is illustrated in Figure 2, which contains a message header and a message body consisting of several Type-Length-Values (TLVs). The TLVs encode Information Elements (IEs) to specify the parameters of a WiMAX message. For example, the MS Info IE includes parameters such as the MAC address and the authenticator ID of the MS. The WiMAX message header consists of the following fields:

- The version field (8 bits) indicates the protocol version number. The current version number is 1.
- The flag field (8 bits) consists of two flags represented by bits. Bit 8 is set if the next expected transaction ID is reset. Bit 7 is set if the source and the destination identifier TLVs are included in the message.
- The function type field (8 bits) indicates the function type (see Table I).
- The OpID field (3 bits) indicates the operation type, which can be 001 (request/initiation), 002 (response), 003 (Ack) or 004 (indication).
- The message type field (5 bits) indicates the message type (see Table I). A function type and a message type uniquely determine a WiMAX message. For example, function type 9 and message type 1 represent the MS_PreAttachment_Req message (this message informs the ASN-GW that an MS attempts to enter the WiMAX network).
- The length field (16 bits) indicates the size of the message in bytes including the header.
- The reserved fields are set to 0.
- The transaction ID field (16 bits) identifies the messages that are delivered in the same transaction.

The TLV format is shown in Figure 3. The type field is a numeric code that represents an IE. The length field specifies the size of the value portion in octets. The value field contains data for this IE (e.g., the MS MAC address for the MS Info IE). The value field itself may contain nested TLVs. Therefore, a WiMAX message

<table>
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<th>Message type</th>
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<td>2 (AuthRelay_EAP_Transfer)</td>
</tr>
<tr>
<td>9 (MS state)</td>
<td>1 (MS_PreAttachment_Req)</td>
</tr>
<tr>
<td></td>
<td>2 (MS_PreAttachment_RSP)</td>
</tr>
<tr>
<td></td>
<td>3 (MS_PreAttachment_Ack)</td>
</tr>
<tr>
<td></td>
<td>4 (MS_Attachment_Req)</td>
</tr>
<tr>
<td></td>
<td>5 (MS_Attachment_RSP)</td>
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<tr>
<td></td>
<td>6 (MS_Attachment_ACK)</td>
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<tr>
<td></td>
<td>7 (Key_Change_Directive)</td>
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</table>

Fig. 2. WiMAX message format.
can be represented as a regular expression [9]:

\[
W = HS^+ , \quad S = TSLV | TNLS^+ \quad (1)
\]

where \( W \) denotes a WiMAX message, \( H \) and \( B \) denote WiMAX message header and body, \( T_S \) and \( T_N \) denote the types for single and nested IE, \( L \) denotes the length for this IE and \( V \) denotes the data for a single IE. This regular expression will be used to decode the WiMAX message in Section 4.

The WiMAX entities interact through a client–server model, where a WiMAX entity plays both server and client roles. We use the initial network entry procedure as an example to illustrate how the WiMAX messages are exchanged between the BS and the ASN-GW. This procedure is executed when an MS enters the WiMAX network [10]. Figure 4 shows the message flow, and the details are given below.

**Step 1.** When an MS enters the WiMAX network, the BS sends an MS_PreAttachment_Req message to the ASN-GW. In this message, the function type is 9, the message type is 1 and the OpID is 001. The flags are set to 0 since the next transaction ID will not be reset and the source/destination identifier TLVs are not included in this message. The IEs of the message are the MS Info (mandatory) and the BS Info (optional).

**Step 2.** The authenticator in the ASN-GW creates a new context block related to this MS and responds to the BS with an MS_PreAttachment_Rsp message. In this message, the function type is 9, the message type is 2, the OpID is 002 and the flags are set to 0. The IEs are the MS Info and the BS Info.

**Step 3.** The BS sends an MS_PreAttachment_Ack message to the authenticator to confirm that the BS has started the negotiation with the MS (the MS is not shown in this figure). In this message, the function type is 9, the message type is 3, the OpID is 003, the bit 8 of the flag is set to 1 and the bit 7 of the flag is set to 0. There is no IE in this message. Note that Steps 1–3 are a transaction initiated by the BS, where the BS is the client and the ASN-GW is the server.

**Step 4.** The extensible authentication protocol (EAP) authentication process is performed between the AAA server, the authenticator in the ASN-GW and the MS (via the BS).

**Step 5.** The authenticator sends a Key_Change_Directive message to the BS to indicate the completion of the EAP authentication process.

**Step 6.** The BS acknowledges with a Key_Change_Ack message, and performs mutual authentication with the MS (not shown in this message flow) [11]. Note that Steps 5 and 6 are a transaction initiated by the ASN-GW, where the ASN-GW is the client and the BS is the server.

**Step 7.** The BS and the authenticator perform the MS attachment to complete the MS registration. Then the ASN-GW establishes the data path.

In this paper, we will show how the above WiMAX procedure can be tested in ITRI-WiMArXT.

### 3. TTCN-3 Test System

The TTCN-3 test system architecture is illustrated in Figure 5, where the system is conceptually divided into several entities that interact with each other during the test execution. This section briefly describes the TTCN-3 entities and interfaces related to ITRI-WiMArXT. Details of the abstract TTCN-3 architecture can be found in References [6,7,12].

#### 3.1. TTCN-3 Test System Entities

The test management (TM; Figure 5 (a)) is responsible for overall management of the system. When the system is initialized, test execution starts with the TM. This entity interacts with the test system user (Figure 5 (i)), and provides test event logs to the user. The TM is also responsible for the invocation of TTCN-3 modules that define an executable test suite. The TM propagates the
module parameters such as implementation extra information for testing (e.g., the IP address of the ASN-GW and the timer values) to the TTCN-3 executable (TE; see Figure 5 (c)). The coding and decoding (CD; see Figure 5 (b)) is responsible for encoding and decoding the data associated with message- or procedure-based communications. The codecs in the CD are implemented in JAVA or C. Therefore, they can be easily ported to different test systems.

The TE is responsible for the interpretation and execution of the test system modules. In the TE, the main test component (MTC; Figure 5 (1)) is responsible for the testcase execution. During the execution, the MTC may create the parallel test components (PTCs; see Figure 5 (2)), and control these PTCs through the test component commands, e.g., create, start and stop (Figure 5 (3)). Each of the PTCs handle one or more interfaces with the system under test (SUT; see Figure 5 (j)). The PTCs are executed concurrently. For example, the TTCN-3 test system may play both roles as the BS and the target ASN-GW to interact with the serving ASN-GW under the test (the SUT) [10]. In this case, the BS and the target ASN-GW are run on two individual PTCs, and the messages exchanged through different interfaces towards the serving ASN-GW are mapped and processed by the corresponding PTCs. The MTC collects the PTCs’ verdicts and determines a global, final verdict for the whole test. Testcase execution ends when the MTC terminates.

The SUT adapter (SA; Figure 5 (d)) bridges the test components and the SUT through the test system interface (TSI; Figure 5 (h)). The SA is implemented in JAVA or C. It propagates the test component’s requests, and then notifies the test component of any response messages from the SUT. The platform adaptor (PA; Figure 5 (e)) adapts the TE to a particular execution platform. It integrates external functionality such as ciphering function and provides the realization of timers (i.e., to start/stop a timer, and to notify the TE when a timeout occurs).

### 3.2. TTCN-3 Test System Interfaces

The **TTCN-3 control interface** (TCI; Figure 5 (f)) [13] and the **TTCN-3 runtime interface** (TRI; Figure 5(g)) [14] are defined as sets of interface operations. These interface operations implement the TTCN-3 operations defined in References [12] and [15]. For example, the TTCN-3 send communication operation invokes the triSend interface operation to send a message to the SUT through a port. Ports provide connection between a test component and the SUT. Some of the TTCN-3 communication operations are listed in Table II. The TM/CD interact with the TE through the TCI. Some of the TCI interface operations invoked by the TE are described as follows:

- The tciGetModulePar operation accesses the values of the module parameters.
- The encode operation encodes a TTCN-3 value into, e.g., a WiMAX message to the SUT.
- The decode operation decodes a message into a TTCN-3 value.

The TRI defines a set of SA interface operations. Some of them invoked by the TE are described as follows:

- The triExecuteTestcase operation is called immediately before a testcase execution. The operation is used to set up, e.g., the parameters of the testcase.
- The triMap operation is invoked by the TTCN-3 map operation to establish a connection to the SUT.
- The triUnmap operation is invoked by the TTCN-3 unmap operation to close the SUT connection, which will release the network resources.
Fig. 6. The mtc\_InitialNetworkEntry testcase.

(1) **testcase mtc\_InitialNetworkEntry () runs on** WimaxMtc **system** WimaxInterface

```cpp
var WimaxBS bs;
bs := WimaxBS.create;
map(bs:client, system:bs.client);
map(bs:server, system:bs.server);
bs.start(tcc\_initialNetworkEntry());
timer t\_initialNetworkEntry;
t\_initialNetworkEntry.start(1.0);
alt{
    [] all component.done {
        t\_initialNetworkEntry.stop;
    };
    [] t\_initialNetworkEntry.timeout {
        all component.stop;
        setverdict(inconc);
    };
}
unmap(bs:client, system:bs.client);
unmap(bs:server, system:bs.server);
}
type component WimaxMtc {}
type component WimaxInterface {}
port WimaxPort bs\_client;
port WimaxPort bs\_server;
... //ports for the target ASN-GW, DHCP server, and so on
}
type component WimaxBS {}
port WimaxPort client;
port WimaxPort server;
}
```

Fig. 6. The mtc\_InitialNetworkEntry testcase.

- The triSend operation is invoked by the TTCN-3 send operation to send an encoded message to the SUT.

An example of the TRI interface operations invoked by the SA is the triEnqueueMsg operation that buffers a message from the SUT. It is called immediately after the SA has received a message.

4. **ITRI-WiMAXT Test System**

Based on the TTCN-3 platform described in the previous section, this section shows how the ITRI-WiMAXT implements TE, TM, CD and SA to test the WiMAX-specific behaviour. We use the mtc\_InitialNetworkEntry testcase in Figure 6 to illustrate how ITRI-WiMAXT works. This test case verifies the initial network entry procedure illustrated in Figure 4. Figure 7 shows the graphical test log for initial network entry procedure and the interaction between the MTC (Figure 7 (1)), PTC (component 1; Figure 7 (2)) and the SUT (system; Figure 7 (3)). Each match box in Figure 7 indicates that the received message matches the pass criteria. The final pass box indicates that the SUT passes this test case.

4.1. **ITRI-WiMAXT Test Executable**

Figure 8 shows the relationship among the components in ITRI-WiMAXT using the unified modelling language [16]. In Figure 6, the test system plays the role of a BS, and the SUT is the serving ASN-GW. The MTC executes on a test component of type WimaxMtc.
The WimaxInterface includes all ports that are used to interact with the SUT in all WiMAX related testcases. For example, the BS PTC ports are defined in Lines 14 and 15 in Figure 6. The MTC may create one or more PTCs (WimaxPtc; see Figure 8 (2)). A WimaxPtc can be WimaxBS (for BS), WimaxASN (for ASN-GW), WimaxDHCP (for DHCP server) and so on. In the mtc_initialNetworkEntry testcase, a WimaxBS called bs (Line 2, Figure 6) is created. This PTC plays the role as a WiMAX BS, and has two ports for the server and the client modes, respectively (see Lines 16 and 17, Figure 6). In the server mode, bs listens on the port server (i.e., WiMAX port 2231/UDP) to wait for transactions initiated from the ASN-GW (the SUT).
In the client mode, \( bs \) initiates a transaction with the port \( client \), which connects to the WiMAX port of the ASN-GW (Lines 3 and 4, Figure 6). Then \( bs \) executes the designated TestCaseControl (Figure 8 (3); see also Line 5, Figure 6). The MTC will wait until either the executions of all PTCs are completed (Line 7, Figure 6), or the timer \( t_{\text{initialNetworkEntry}} \) is expired (Lines 8–10, Figure 6). When the testcase is complete, the \texttt{unmap} operation is invoked to release the ports (Lines 11 and 12, Figure 6).

To process the WiMAX messages, the ITRI-WiMAXT defines a TTCN-3 record type \texttt{WimaxMsg} (Figure 8 (5)) based on the WiMAX specification described in Figure 2. The WiMAX message header and the message body (i.e., the IEs) are implemented as the \texttt{WimaxMsgHeader} and the \texttt{IEList} (Lines 1 and 2, Figure 9). The \texttt{IEList} consists of a set of \texttt{InfoElements} (Line 3, Figure 9). At Line 4 of Figure 9, the \texttt{InfoElement} is a TTCN-3 record type representing the IE of the WiMAX message described in Figure 3. The \texttt{InfoElement} contains an integer \texttt{ieType} and an \texttt{IEValue} type \texttt{ieValue} (Lines 5 and 6, Figure 9), which represent the type and the value fields of the IE, respectively. The \texttt{IEValue} is a TTCN-3 union type, which can be single or nested. A single \texttt{IEValue} is a hex string (Line 9, Figure 9). A nested \texttt{IEValue} includes another \texttt{IEList} (Line 8, Figure 9). Note that the length fields of the WiMAX message and the IE are calculated in the CD, which are not presented in the \texttt{WimaxMsg} and the \texttt{InfoElement}.

Parameter values of a specific \texttt{WimaxMsg} are filled in the \texttt{WimaxTemplate} (Figure 8 (4)). Figure 10 shows how the \texttt{MS_PreAttachment_REQ} message is parameterized in the \texttt{WimaxTemplate} called \texttt{msPreAttachmentReq}. According to the description in Section 2,
modulepar {
(1) charstring ASNGW.IP ADDRESS := "140.113.214.113";
(2) charstring BS.IP ADDRESS := "140.113.214.114";
(3) hexstring MSID := '001123344556H;
... // other WiMAX parameters
}

Fig. 12. Examples of the module parameters.

the values of the version, flag, funcType, opID and msgType fields for this message are set to 1, 0, 9, 1 and 6, respectively (Lines 2–6 in Figure 10). Fields MSID, transaction ID and TLVs are determined during the test-case execution, and are passed to the WimaxTemplate through the template parameters _a_msId, _a_transId and _a_ieList (Lines 1, 7–9 in Figure 10).

A WiMAX procedure is executed by the TestCaseControl through sending and receiving a sequence of WimaxTemplates. For example, the tcc_initialNetworkEntry TestCaseControl for mtc_initialNetworkEntry is illustrated in Figure 11. Lines 1–3 in Figure 11 correspond to Steps 1–3 in Figure 4, where the BS is in the client mode. Specifically, the TestCase sends an MS_PreAttachment_REQ message (through WimaxTemplate msPreAttachmentReq in Figure 10) to the SUT at Line 1, where client.send invokes f_getTransId in the FunctionLibrary (Figure 8

```java
public TriMessage encode(final Value TTCNValue) {
    byte[] encMsg = new byte[65535];
    (1) RecordValue rv = (RecordValue) TTCNValue;
    (2) if(rv.getType().getName().equals("WimaxMsg")) { // is it a WimaxMsg?
        int msgLen;
        try {
            (3) encMsg[0] = (byte)((IntegerValue)rv.getField("version").getInt());
            (4) ... // encodes the header
                List _ieList = (List)rv.getField("ieList");
            (5) if( !_ieList.isEmpty() ) { // does the message has the body?
                msgLen = encodeIEList(encMsg, _ieList, 20);
            }
            (7) encMsg[4] = (byte) (msgLen/256);
            (8) encMsg[5] = (byte) (msgLen%256);
        } catch (Exception e) {
            e.printStackTrace();
        }
    } else ... // encode other message types such as DHCP
    (12) ... //trim the encMsg
    (13) return = new TriMessageImpl(encMsg);
}

private int encodeIEList(byte[] encMsg, List _ieList, int ptr){
    int plen = _ieList.getLength();
    (14) for(int i = 0; i < plen; i++){
        (15) Record ie = (Record)_ieList.getDataItem(i);
        (16) int T = ((IntegerValue) ie.getField("ieType").getInt());
        (17) encMsg[ptr++] = (byte) (T / 256);
        (18) encMsg[ptr++] = (byte) (T % 256);
        (19) ... // encode the value, and then calculate the length
    }
    return ptr;
}
```

Fig. 13. ITRI-WiMAXT encode operation.
(6)) to generate a new transaction ID. Line 2 of Figure 11 checks whether the received msPreAttachmentRsp is correct. Then Line 3 sends msPreAttachmentAck to the SUT (ASN-GW). Lines 4 and 5 in Figure 11 correspond to Steps 5 and 6 in Figure 4, where the BS is in the server mode. If all messages from the ASN-GW are successfully verified, the verdict of this testcase is set to pass. Otherwise, the verdict is set to fail.

4.2. ITRI-WiMAXT TM and CD Entities

The TM and the CD entities interact with the TE through the TCI interface. In ITRI-WiMAXT, some module parameters passed from the TM to the TE are illustrated in Figure 12, which include the IP addresses for the ASN-GW and the BS (Lines 1 and 2 in Figure 12), and an MSID (Line 3 in Reference 12). These module parameters can be dynamically set by the test system user, and are retrieved by the triExecuteTestcase operation for the testcase execution.

As implemented in JAVA, the ITRI-WiMAXT CD encodes and decodes the WiMAX messages. When client.send in Figure 11 is executed, the TE invokes the encode operation in the CD, which encodes a TTCN value TTCNValue (e.g., msPreAttachmentReq in Figure 10) into an encoded byte array encMsg. Line 1 in Figure 13 stores TTCNValue into a variable rv. Line 2 checks if a WiMAX message should be created. If so, Lines 3–10 encode rv based on the WiMAX format. Specifically, Lines 3 and 4 in Figure 13 encode the WiMAX message

```java
public Value decode(final TriMessage encTriMsg, final Type decodingHypothesis) {
    if (decodingHypothesis.getName().equals("WimaxMsg")) { // is it a WimaxMsg
        final RecordValue TTCNValue = (RecordValue) decodingHypothesis.newInstance();
        encMsg = encTriMsg.getMessage();
        return createWimaxMsg(TTCNValue, encMsg);
    } else ... // decode other message types such as DHCPMsg
}

private RecordValue createWimaxMsg(final RecordValue TTCNValue, final byte[] encMsg) {
    ... // decode the message header from the first 20 bytes
    if (hasIEList(TTCNValue)) { // decode the message body
        Value value = decodeIEList((RecordOfValue) TTCNValue.getField("ieList"),
                                   encMsg, 20, encMsg.length);
        setField(TTCNValue, "ieList", value);
    }
    return TTCNValue;
}

private Value decodeIEList(RecordOfValue ieList, final byte[] encMsg, int ptr, int lptr) {
    ieList.setLength(0);
    while (ptr != lptr) {
        ... // decode and obtain the type T and length L of the IE from the first four bytes
        if (matches(T, T5)) { // is a single IE type matched?
            ... // obtain the value V
        } else { // is a nested IE type matched?
            ... // recursively decode the information element
            V = decodeIEList(...);
        }
        ... // insert the T, L and V as an IE into the ieList
    }
    return ieList;
}
```

Fig. 14. ITRI-WiMAXT decode operation.
header of \( rv \). Line 5 checks whether the WiMAX message has a body. If so, Line 6 encodes the message body by calling the function \texttt{encodeIEList} with parameters such as \texttt{encMsg}, \texttt{ieList} and the encoding position \( ptr \). This operation repeatedly encodes the IEs (Lines 14–19). In each loop, an IE \( ie \) is retrieved from \texttt{ieList} (Line 15). The operation encodes the type and value fields, and then calculate the length of the IE (Lines 16–19). The encoded results are saved in \texttt{encMsg}. Then \( ptr \) is returned to \texttt{encode} for calculating the message length \( msgLen \). Lines 7 and 8 in Figure 13 encode \texttt{msgLen} and fill it into \texttt{encMsg}. If an exception occurs (e.g., the array has been accessed with an illegal index), the debug information is printed (Lines 9 and 10). Lines 12 and 13 trim the unused tailing spaces of the byte array and return it.

When \texttt{client.receive} in Figure 11 is executed, the CD \texttt{decode} operation in Figure 14 is invoked, where an encoded TriMessage \texttt{encTriMsg} is decoded based on the \texttt{decodingHypothesis} type that can be WimaxMsg, DHCPMsg and so on. At Lines 1 and 2 in Figure 14, if \texttt{decodingHypothesis} indicates the type ‘WimaxMsg’, the \texttt{newInstance} operation creates an empty TTCN value \texttt{TTCNValue}. Line 3 retrieves the encoded byte array \texttt{encMsg} from \texttt{encTriMsg}. In Line 4, the \texttt{createWimaxMsg} operation is called to fill the fields in \texttt{TTCNValue}. Specifically, this operation decodes the WiMAX message header from the first 20 bytes of the \texttt{encMsg} (Line 5), and then decodes the remaining bytes as the WiMAX IEs (Lines 6–9) by calling the \texttt{decodeIEList} operation with parameters such as \texttt{ieList} (the WiMAX message body at Line 2 in Figure 9), the \texttt{encMsg} and the range of the remaining bytes (i.e., from currently decoding position to the end). This operation uses the regular expression (1) to repeatedly decode the IEs (Lines 10–16) until all bytes in \texttt{encMsg} are traversed. In each loop, the type field \( T \) and the length field \( L \) are decoded from the first 4 bytes (Line 11). If \( T \) matches the \( T_3 \) type described in Section 2 (i.e., it is a single IE), the \( L \)-bytes value field is filled (Lines 12 and 13). Otherwise, it is a nested IE, and is recursively decoded into a value \( V \) (Lines 14 and 15). Finally, the \( T \), \( L \) and \( V \) are inserted into the \texttt{ieList} as an IE (Line 16). When all bytes are decoded, the value \texttt{ieList} is returned (Line 17).

5. Conclusions

This paper described the architecture and implementation of ITRI-WiMAXT, a WiMAX test system developed based on the TTCN-3 specifications. This system has been jointly developed by the National Telecommunication Program (NTP) and the Industrial Technology Research Institute (ITRI) in Taiwan. We showed how to encode and decode the WiMAX messages in ITRI-WiMAXT, and used the initial network entry procedure as an example to illustrate how conformance tests can be implemented and conducted for WiMAX. A major advantage of ITRI-WiMAXT is its modular implementation. The implemented CDs can be reused in the new testcases.

References

Authors’ Biographies

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